# Digital Geometry Processing 

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x-ray diffractometer
Geometry
$\gamma \varepsilon \omega \mu \varepsilon \tau \rho i \alpha$


## geo $=$ earth metria $=$ measure


radio telescope

laser scanner

## Digital Geometry

- Entertainment Industry

- Modeling $\rightarrow$ digital character \& set design
- Simulation $\rightarrow$ computer games, movies, special effects


## Digital Geometry

- Medical Applications


MRI scanner


2D slices

- Analysis $\rightarrow$ diagnosis, operation planning
- Modeling $\rightarrow$ design of prosthetics
- Simulation $\rightarrow$ surgery training


## Digital Geometry

- Engineering Applications

- Analysis $\rightarrow$ quality control
- Modeling $\rightarrow$ product design, rapid prototyping
- Simulation $\rightarrow$ aerodynamics, crash tests


## Digital Geometry

- 3D City Modeling


range-data, images, etc.


3D city model

- Analysis $\rightarrow$ navigation, map design
- Modeling $\rightarrow$ urban planning, virtual worlds
- Simulation $\rightarrow$ traffic, pollution, etc.


## Application Areas

- Computer games
- Movie production
- Engineering
- Cultural Heritage
- Topography
- Architecture
- Medicine
- etc.



## Geometry Processing Pipeline



## Geometry Processing Pipeline



## Geometry Processing Pipeline



Surface smoothing for noise removal


## Geometry Processing Pipeline



Surface smoothing for noise removal

$\downarrow$


## Geometry Processing Pipeline



## Geometry Processing Pipeline



Simplification for complexity reduction


Remeshing for improving mesh quality


## Geometry Processing Pipeline



Freeform and multiresolution modeling


## Geometry Processing Toolbox

- Geometric Modeling
- Methods \& algorithms for representing and processing geometric objects
- Geometry processing
- Core algorithms?
- Efficient implementations?


# Shape Reconstruction 

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## Outline

- Sensors
- Problem statement
- Computational Geometry
- Voronoi/Delaunay
- Alpha-shapes
- Crust
- Variational formulations
- Poisson reconstruction

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## SENSORS

## Laser scanning


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## Car-based Laser



## Airborne Lidar



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## Multi-View Stereo (MVS)



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## Depth Sensors



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PROBLEM STATEMENT

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## Reconstruction Problem

Input: point set $P$ sampled over a surface $S$ :

Non-uniform sampling
With holes
With uncertainty (noise)

point set

Output: surface
Approximation of $S$ in terms of topology and geometry

Desired:
Watertight
Intersection free

reconstruction

surface

## Ill-posed Problem



Many candidate surfaces for the reconstruction problem!

## III-posed Problem



Many candidate surfaces for the reconstruction problem! How to pick?

Priors


Smooth


Piecewise Smooth

"Simple"

## Surface Smoothness Priors




Global: linear, eigen, graph cut, ...
Robustness to missing data

Piecewise Smoothness


Sharp near features
Smooth away from features

## Domain-Specific Priors



## Previous Work









































 Computers.and Graphics, 11:393\{408\}, 1987,

## Warm-up



Smooth


Piecewise Smooth

"Simple"
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VORONOI / DELAUNAY

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## Voronoi Diagram

Let $\mathcal{E}=\left\{\mathbf{p}_{\mathbf{1}}, \ldots, \mathbf{p}_{\mathbf{n}}\right\}$ be a set of points (so-called sites) in $\mathbb{R}^{d}$. We associate to each site $\mathbf{p}_{\mathbf{i}}$ its Voronoi region $V\left(\mathbf{p}_{\mathbf{i}}\right)$ such that:

$$
V\left(\mathbf{p}_{\mathbf{i}}\right)=\left\{\mathbf{x} \in \mathbb{R}^{d}:\left\|\mathbf{x}-\mathbf{p}_{\mathbf{i}}\right\| \leq\left\|\mathbf{x}-\mathbf{p}_{\mathbf{j}}\right\|, \forall j \leq n\right\} .
$$


http://www.cgal.org

## Delaunay Triangulation

## Dual structure of the Voronoi diagram.

The Delaunay triangulation of a set of sites E is a simplicial complex such that $k+1$ points in $E$ form a Delaunay simplex if their Voronoi cells have nonempty intersection


## Empty Circle Property

Empty circle: A triangulation $T$ of a point set $E$ such that any d-simplex of $T$ has a circumsphere that does not enclose any point of E is a Delaunay triangulation of E . Conversely, any $k$-simplex with vertices in E that can be circumscribed by a hypersphere that does not enclose any point of $E$ is a face of the Delaunay triangulation of $E$.


## Delaunay-based

Key idea: assuming dense enough sampling, reconstructed triangles are Delaunay triangles.


## Alpha-Shapes [Edelsbrunner, Kirkpatrick, Seidel]



Segments: point pairs that can be touched by an empty disc of radius alpha.

## Alpha-Shapes

In 2D: family of piecewise linear simple curves constructed from a point set $P$.

Subcomplex of the Delaunay triangulation of $P$.
Generalization of the concept of the convex hull.


## Alpha-Shapes


$\alpha=0 \quad$ Alpha controls the desired level of detail.

$\alpha=\infty$

## Delaunay-based

Key idea: assuming dense enough sampling, reconstructed triangles are Delaunay triangles.

First define
Medial axis
Local feature size
Epsilon-sampling


MEDIAL AXIS

Ennis

## Medial Axis

For a shape (curve/surface) a Medial Ball is a circle/sphere that only meets the shape tangentially, in at least two points.


## Medial Axis

For a shape (curve/surface) a Medial Ball is a circle/sphere that only meets the shape tangentially, in at least two points.
The centers of all such balls make up the medial axis/skeleton.


## Medial Axis



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## Medial Axis

## Observation*:

For a reasonable point sample, the medial axis is wellsampled by the Voronoi vertices.
*In 3D, this is only true for a subset of the Voronoi vertices - the poles.

## Voronoi \& Medial Axis



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## Local Feature Size



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## Epsilon-Sampling



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## Crust [Amenta et al. 98]

If we consider the Delaunay Triangulation of a point set, the shape boundary can be described as a subset of the Delaunay edges.

Q: How do we determine which edges to keep?
A: Two types of edges:

1. Those connecting adjacent points on the boundary
2. Those traversing the shape.

Discard those that traverse.


## Crust [Amenta et al. 98]

## Observation:

Edges that traverse cross the medial axis.
Although we don't know the axis, we can sample it with the Voronoi vertices.

Edges that traverse must be near the Voronoi vertices.


## Crust [Amenta et al. 98]



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## Delaunay Triangulation



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Delaunay Triangulation \& Voronoi Diagram


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Voronoi Vertices

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## Refined Delaunay Triangulation



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Crust


## Crust



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## Crust (variant)

Algorithm:

1. Compute the Delaunay triangulation.
2. Compute the Voronoi vertices
3. Keep all edges for which there is a circle that contains the edge but no Voronoi vertices.


## SPECTRAL «CRUST »

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## Space Partitioning

Given a set of points, construct the Delaunay triangulation.

If we label each triangle as inside/outside, then the surface of interest is the set of edges that lie between inside and outside triangles.


## Space Partitioning

Q: How to assign labels?
A: Spectral Partitioning
Assign a weight to each edge indicating if the two triangles are likely to have the same label.


## Space Partitioning

## Assigning edge weights

Q: When are triangles on opposite sides of an edge likely to have the same label?

A: If the triangles are on the same side, their circumscribing circles intersect deeply.
Use the angle of intersection to set the weight.


## Crust

## Several Delaunay algorithms provably correct



## Delaunay-based

Several Delaunay algorithms are provably correct... in the absence of noise and undersampling.
perfect data?

Noise \& Undersampling


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## Delaunay-based

Several Delaunay algorithms are provably correct... in the absence of noise and undersampling.

Motivates reconstruction by fitting approximating implicit surfaces


## VARIATIONAL FORMULATIONS



Smooth


Piecewise Smooth

"Simple"

## Poisson Surface Reconstruction

[Kazhdan et al. 06]

## Indicator Function

Construct indicator function from point samples


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## Indicator Function

Construct indicator function from point samples


## Poisson Surface Reconstruction



Mana

## Poisson Surface Reconstruction



WHAT NEXT

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## What Next

Online

- Reconstruction
- Localization

Robustness

- Structured outliers
- Heterogeneous data


## «La Lune»



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## « La Lune »


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# Simplification \& Approximation 

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## Outline

- Motivations
- Simplification
- Approximation
- Remaining Challenges


## Motivations

- Multi-resolution hierarchies for
- efficient geometry processing
- level-of-detail (LOD) rendering



## Complexity-Error Tradeoff



## Problem Statement

- Given: $\mathcal{M}=(\mathcal{V}, \mathcal{F})$
- Find: $\mathcal{M}^{\prime}=\left(\mathcal{V}^{\prime}, \mathcal{F}^{\prime}\right)$ such that

1. $\left|\mathcal{V}^{\prime}\right|=n<|\mathcal{V}|$ and $\left\|\mathcal{M}-\mathcal{M}^{\prime}\right\|$ is minimal, or
2. $\left\|\mathcal{M}-\mathcal{M}^{\prime}\right\|<\epsilon$ and $\left|\mathcal{V}^{\prime}\right|$ is minimal


## Problem Statement

- Given: $\mathcal{M}=(\mathcal{V}, \mathcal{F})$
- Find: $\mathcal{M}^{\prime}=\left(\mathcal{V}^{\prime}, \mathcal{F}^{\prime}\right)$ such that

1. $\left|\mathcal{V}^{\prime}\right|=n<|\mathcal{V}|$ and $\left\|\mathcal{M}-\mathcal{M}^{\prime}\right\|$ is minimal, or
2. $\left\|\mathcal{M}-\mathcal{M}^{\prime}\right\|<\epsilon$ and $\left|\mathcal{V}^{\prime}\right|$ is minimal
hard! [Agarwal-Suri 1998]
$\rightarrow$ look for sub-optimal solution

## Simplification

## Simplification

- Vertex Clustering
- Iterative Decimation
- Extensions


## Simplification

- Vertex Clustering
- Iterative Decimation
- Extensions


## Vertex Clustering

- Cluster Generation
- Uniform 3D grid
- Map vertices to cluster cells
- Computing a representative
- Mesh generation

- Topology changes


## Vertex Clustering

- Cluster Generation
- Hierarchical approach
- Top-down or bottom-up
- Computing a representative
- Mesh generation
- Topology changes


## Vertex Clustering

- Cluster Generation
- Computing a representative
- Average/median vertex position
- Error quadrics
- Mesh generation
- Topology changes


## Computing a Representative



- Average vertex position $\rightarrow$ Low-pass filter


## Computing a Representative



- Median vertex position $\rightarrow$ Sub-sampling


## Computing a Representative



- Error quadrics


## Error Quadrics

- Squared distance to plane

$$
\begin{gathered}
p=(x, y, z, 1)^{T}, \quad q=(a, b, c, d)^{T} \\
\operatorname{dist}(q, p)^{2}=\left(q^{T} p\right)^{2} \\
Q_{q}=\left[\begin{array}{llll}
a^{2} & a b & a c & a d \\
a b & b^{2} & b c & b d \\
a c & b c & b^{2} & c d \\
a d & b d & c d & d^{2}
\end{array}\right]
\end{gathered}
$$

## Error Quadrics

- Sum distances to vertex' planes
$\sum_{i} \operatorname{dist}(q, p)^{2}$
- Point location that minimizes the error

$$
\left[\begin{array}{cccc}
q_{11} & q_{12} & q_{13} & q_{14} \\
q_{21} & q_{22} & q_{23} & q_{24} \\
q_{31} & q_{32} & q_{33} & q_{34} \\
0 & 0 & 0 & 1
\end{array}\right] p^{*}=\left[\begin{array}{l}
0 \\
0 \\
0 \\
1
\end{array}\right]
$$

## Vertex Clustering

- Cluster Generation
- Computing a representative
- Mesh generation
- Clusters $p\left\{p_{0}, \ldots, p_{n}\right\}, q\left\{q_{0}, \ldots, q_{m}\right\}$
- Connect ( $p, q$ ) if there was an edge ( $p_{i}, q_{j}$ )
- Topology changes


## Vertex Clustering

- Cluster Generation
- Computing a representative
- Mesh generation
- Topology changes
- If different sheets pass through one cell
- Not manifold



## Simplification

- Vertex Clustering
- Iterative Decimation
- Extensions


## Iterative Decimation

- General Setup
- Decimation operators
- Error metrics
- Fairness criteria
- Topology changes


## General Setup

Repeat:

- pick mesh region
- apply decimation operator

Until no further reduction possible

## Greedy Optimization

For each region

- evaluate quality after decimation
- enqueue(quality, region)

Repeat:

- pick best mesh region
- apply decimation operator
- update queue

Until no further reduction possible

## Iterative Decimation

- General Setup
- Decimation operators
- Error metrics
- Fairness criteria
- Topology changes


## Decimation Operators

- What is a "region" ?
- What are the DOF for re-triangulation?
- Classification
- Topology-changing vs. topology-preserving
- Subsampling vs. filtering


## Vertex Removal



Select a vertex to be eliminated

## Vertex Removal



Select all triangles
sharing this vertex

## Vertex Removal



Remove the selected triangles, creating the hole

## Vertex Removal



Fill the hole with triangles

## Decimation Operators



- Remove vertex
- Re-triangulate hole
- Combinatorial DOFs
- Sub-sampling


## Decimation Operators



- Merge two adjacent triangles
- Define new vertex position
- Continuous DOF
- Filtering


## Decimation Operators



- Collapse edge into one end point
- Special vertex removal
- Special edge collapse
- No DOFs
- One operator per half-edge
- Sub-sampling


## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Edge Collapse



## Incremental Decimation

- General Setup
- Decimation operators
- Error metrics
- Fairness criteria
- Topology changes


## Local Error Metrics

- Local distance to mesh [schroeder et al. 92]
- Compute average plane
- No comparison to original geometry



## Local Error Metrics

- Volume preserving [Lindstrom-Turk]. Fast and memory efficient polygonal simplification. IEEE Visualization 98.


Implemented in

## Global Error Metrics

- Simplification envelopes [Cohen et al. 96]
- Compute (non-intersecting) offset surfaces
- Simplification guarantees to stay within bounds



## Global Error Metrics

- (Two-sided) Hausdorff distance: Maximum distance between two shapes
- In general $d(A, B) \neq d(B, A)$
- Compute-intensive


Valette et al. Mesh Simplification using a

## Global Error Metrics

- One-sided Hausdorff distance
- From original vertices to current surface



## Global Error Metrics

- Error quadrics [Garland, Heckbert 97]
- Squared distance to planes at vertex
- No bound on true error



## Error Quadrics



## Incremental Decimation

- General Setup
- Decimation operators
- Error metrics
- Fairness criteria
- Topology changes


## Fairness Criteria

- Rate quality of decimation operation
- Approximation error
- Triangle shape
- Dihedral angles
- Valence balance
- Color differences



## Fairness Criteria

- Rate quality after decimation
- Approximation error
- Triangle shape
- Dihedral angles
- Valance balance
- Color differences



## Incremental Decimation

- General Setup
- Decimation operators
- Error metrics
- Fairness criteria
- Topology changes


## Topology Changes

- Merge vertices across non-edges
- Changes mesh topology
- Need spatial neighborhood information
- Generates non-manifold meshes



## Topology Changes

- Merge vertices across non-edges
- Changes mesh topology
- Need spatial neighborhood information
- Generates non-manifold meshes

manifold

non-manifold

Approximation

## Variational Shape Approximation

- Rationale: cast surface approximation as a variational k-partitioning problem


Cohen-Steiner, A., Desbrun.
Variational Shape Approximation. SIGGRAPH 2004.

## Simpler Setting: 2D Partitioning



## Energy

$$
E=\sum_{j=1 . . k} \int_{x \in R_{j}} \rho(x)\left\|x-x_{j}\right\|^{2} d x
$$


density function

## Lloyd Iteration

- Alternate:
- Voronoi partitioning
- Relocate sites to centroids
- Minimizes energy
- Necessary condition for optimality: Centroidal Voronoi
 tessellation


## Variational Shape Approximation

- Rationale: cast surface approximation as a variational k-partitioning problem
- for each region, find best-fit linear proxy
- "best fit" for a given metric



## Variational Shape Approximation

- Distortion
= integrated error between region and proxy
- Total distortion = sum of proxy distortion
- Best k-approximation = minimum distortion


## Overview


initial mesh

+ partition

associated proxies

proxy-based remeshing


## K-Means Clustering

Starting with k-generators

Alternate:

- cluster by closest proximity (creates regions $R_{j}$ )
- find new generators $c_{j}$ of regions $R_{j}$


## Partition Optimization

Clustering for Approximation

- Replace points by proxies
- Min approximation error
- Equi-distribute energy among proxies


## Error Metrics

- $L^{2}$
- asymptotically, aspect ratio is $\sqrt{\kappa_{1} / \kappa_{2}}$
- hyperbolic regions troublesome
- no unique minimum
- convergence in $L^{2}$ does not guarantee in normals
- example: Schwarz's Chinese lantern
- [Shewchuck 04] gradient bounds harder than interpolation
- $\mathrm{L}^{2,1} \iint_{x \in X}\left\|\mathbf{n}(x)-\mathbf{n}_{i}\right\|^{2} d x$
- asymptotically, aspect ratio is $\kappa_{1} / \kappa_{2}$
- hyperbolic regions ok
- captures normal field
$L^{2}$ VS. $L^{2,1}$



## Triangulation

- node vertex
- where 3+ regions meet
- 2+ on boundary



## Triangulation

- node wedge



## Triangulation

- Two-pass flooding algorithm (~multi-source Djisktra's shortest path algorithm)
- first pass: flood only region boundaries (to enforce the constrained edges)
- second pass: flood interior areas



## First pass



## Second pass



## Triangulation



## Triangulation




## Metrics



## Example



## Example



## Example







Remaining Challenges

## Remaining Challenges

- Beyond approximation
- Abstraction [Sheffer, Mitra et al. 2009] Abstraction of Man-Made Shapes.



## Remaining Challenges

- Beyond approximation
- Meaningful LODs. [Verdié, Lafarge, A. 2013]


